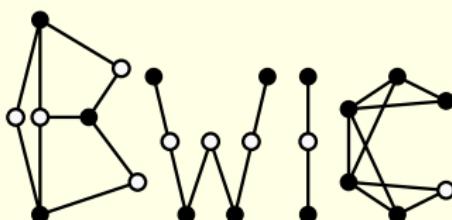


Tolerant identifying codes using Euclidean balls

Aline Parreau

With Ville Junnila and Tero Laihonen



Bordeaux Workshop on Identifying Codes
November 31-25, 2011

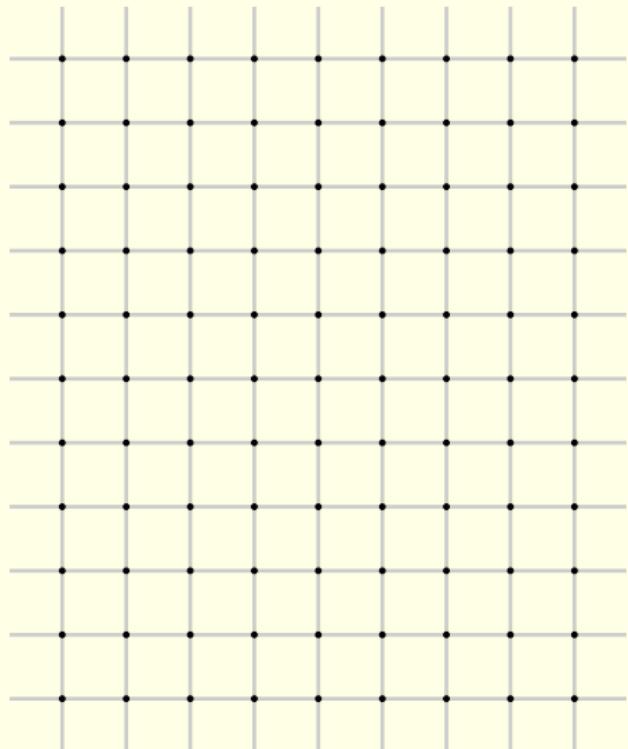
ANR IDEA 

maths à modeler



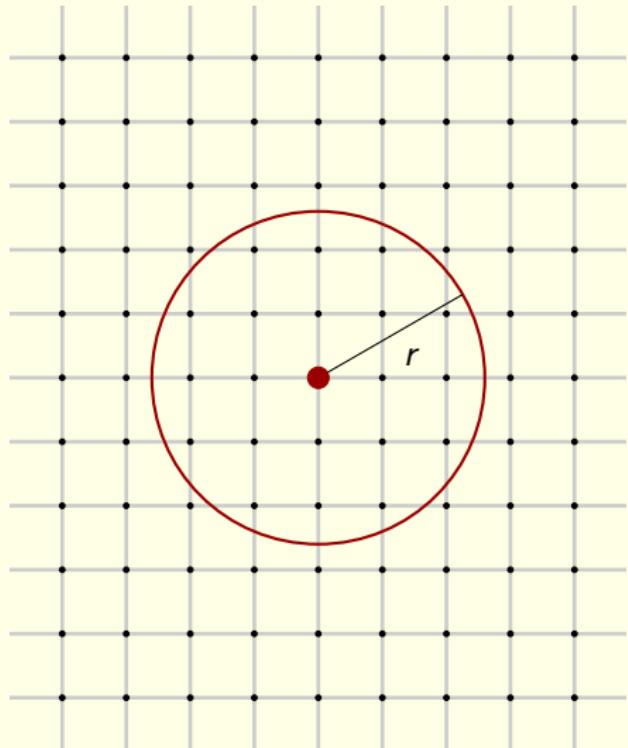
Framework

- Vertices (sensors): \mathbb{Z}^2



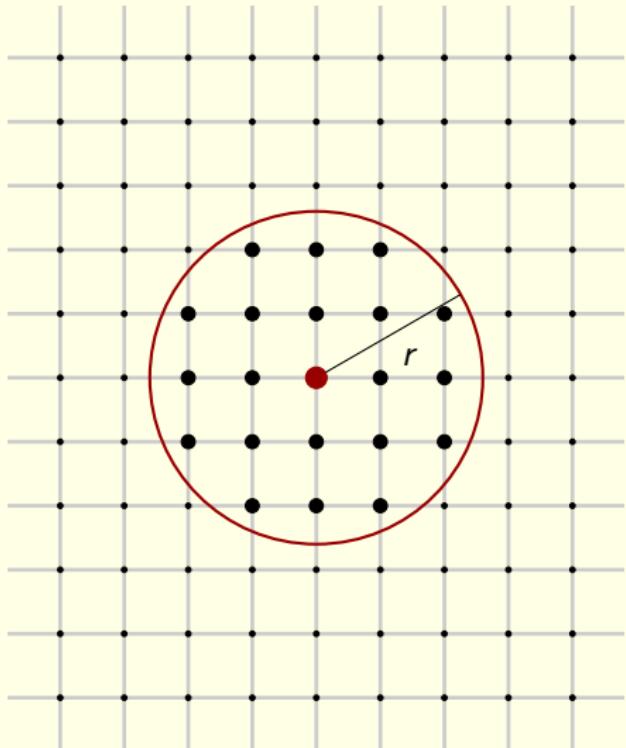
Framework

- Vertices (sensors): \mathbb{Z}^2
- Can detect faults at Euclidean distance r



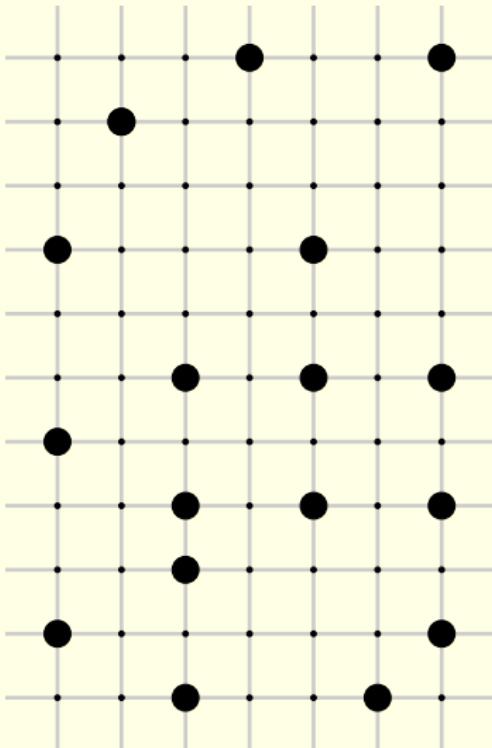
Framework

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→ Euclidean grid



Identifying codes in Euclidean grids

Identifying code : subset $C \subseteq \mathbb{Z}^2$ such that:

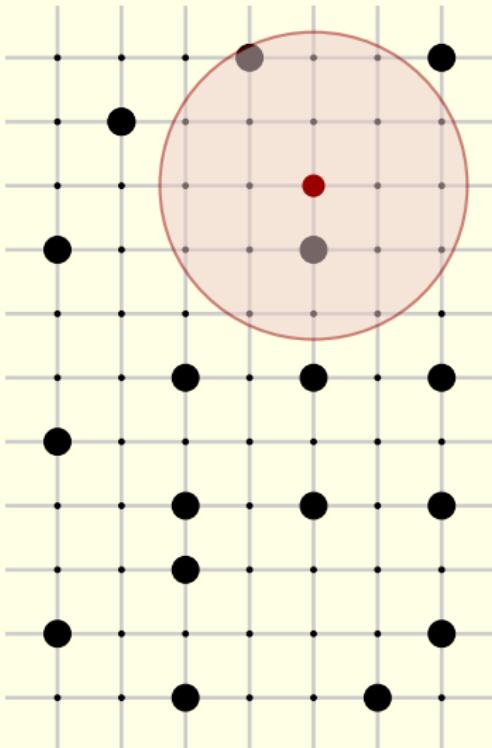


Identifying codes in Euclidean grids

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- C is an r -dominating set: for all $u \in \mathbb{Z}^2$,

$$C \cap B_r(u) \neq \emptyset$$



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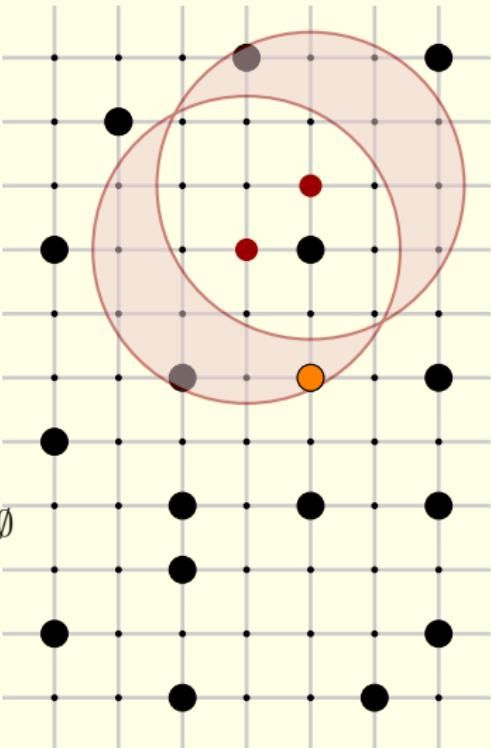
- C is an **r -dominating set**: for all $u \in \mathbb{Z}^2$,

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- C is an **r -separating set** : for all $u, v \in \mathbb{Z}^2$,

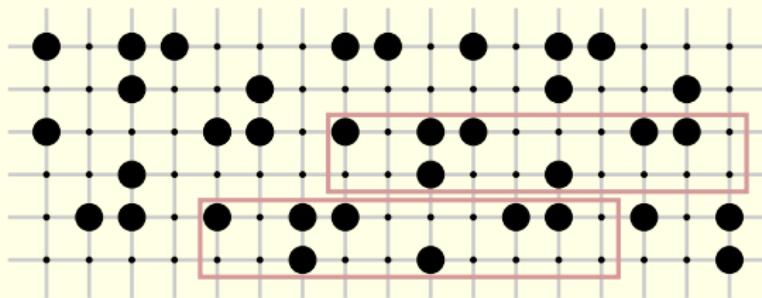
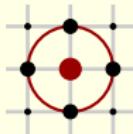
$$C \cap ((B_r(u) \setminus B_r(v)) \cup (B_r(v) \setminus B_r(u))) \neq \emptyset$$

Goal: minimize **density** $D(r)$ of C



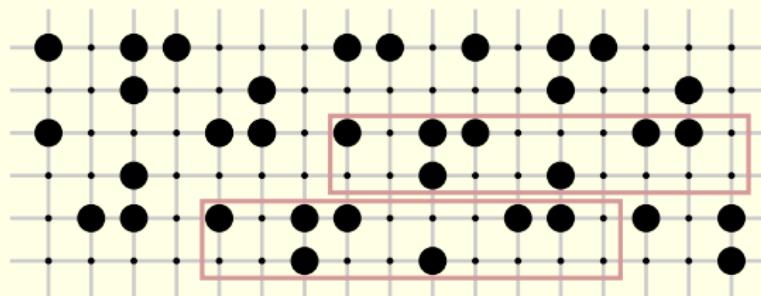
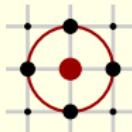
Square grid

$r = 1 \rightarrow$ square grid



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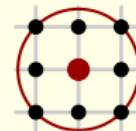
Proposition Cohen et al, 1999 and Ben-Haim & Litsyn, 2005

The optimal density of an identifying code in the square grid is $D(1) = \frac{7}{20}$

Other results for small values of r

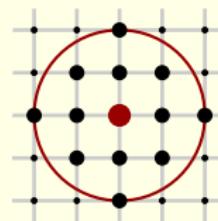
$r = \sqrt{2}$, King grid:

$$D(\sqrt{2}) = \frac{2}{9}$$



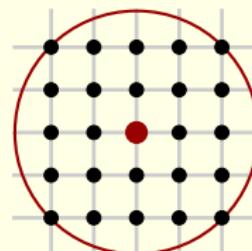
$r = 2$, Square grid of radius 2:

$$\frac{6}{35} \leq D(2) \leq \frac{5}{29}$$



$r = 2\sqrt{2}$ King grid of radius 2:

$$D(2\sqrt{2}) = \frac{1}{8}$$



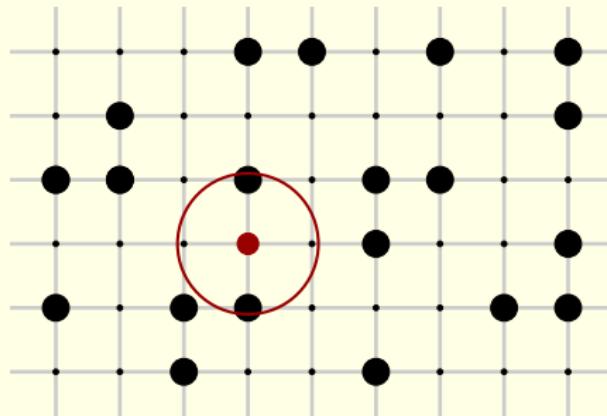
For a general radius

Proposition Junnila and Laihonen, 2011

General bounds:

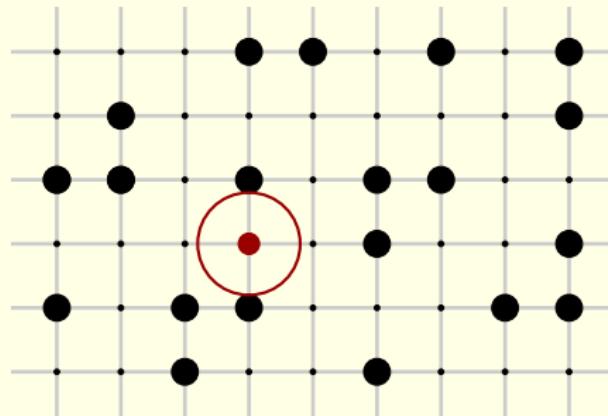
$$\frac{1}{3.22r + 4} \leq D(r) \leq \frac{1}{2[r]}$$

Optimal codes are not robust



The code is not any more identifying if the radius changes a little.

Optimal codes are not robust



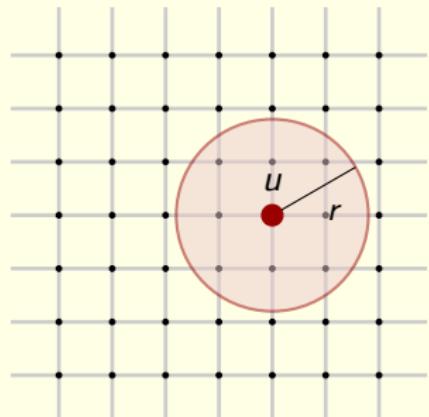
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Tolerant identifying codes

Balls are between balls of radius r and $r + \Delta$.

$C \subseteq \mathbb{Z}^2$ is an (r, Δ) -identifying code
if C is:

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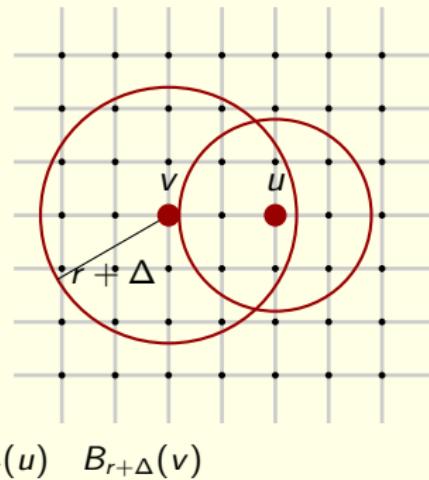
$B_r(u)$

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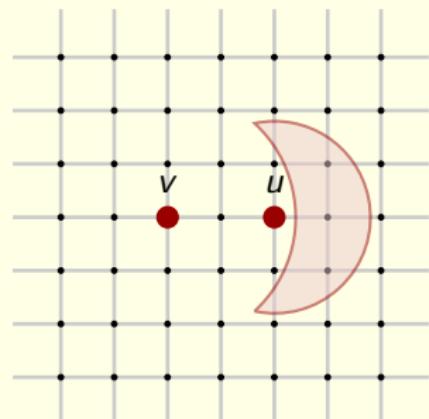
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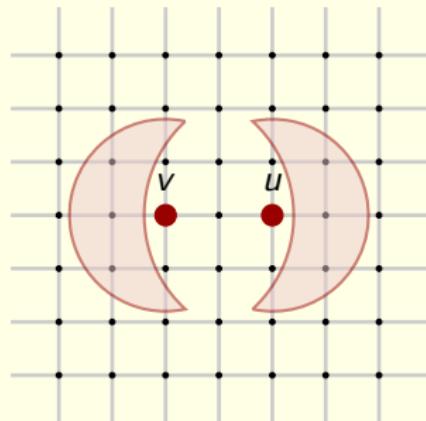
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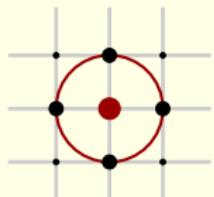
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$D(r, \Delta)$: minimal density

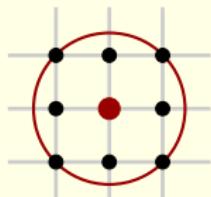


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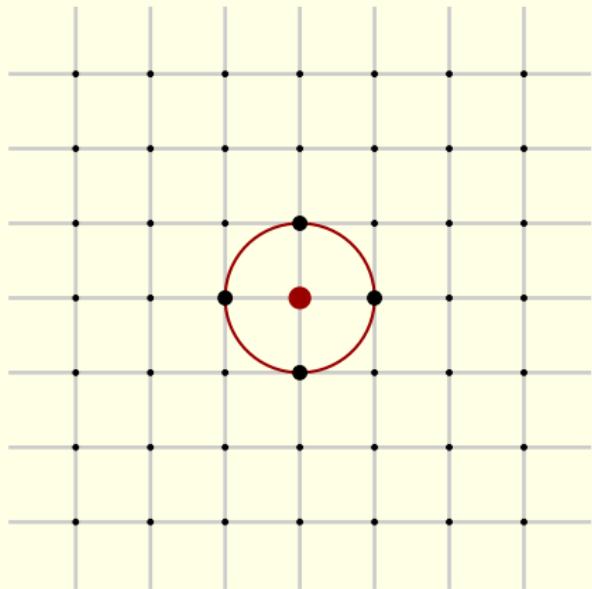
First example: $(r, \Delta) = (1, \sqrt{2} - 1)$



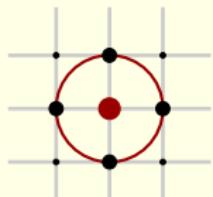
$$r = 1$$



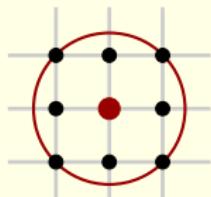
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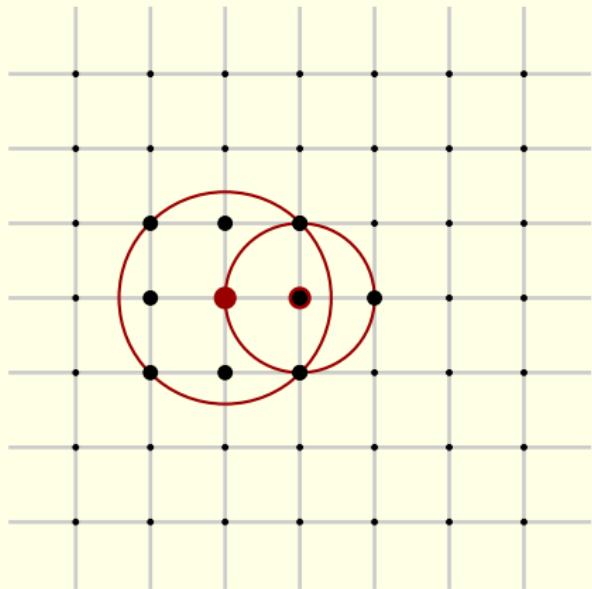
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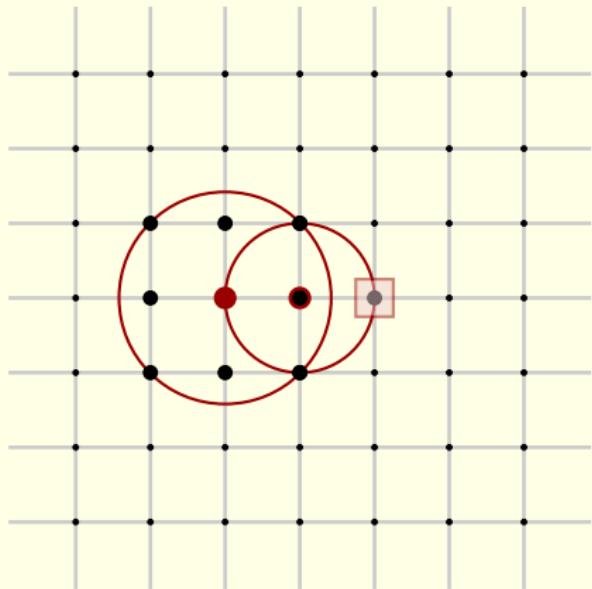
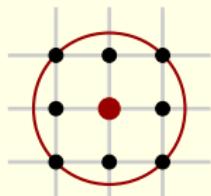
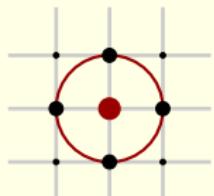
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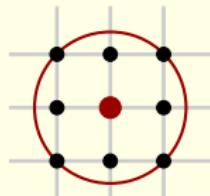
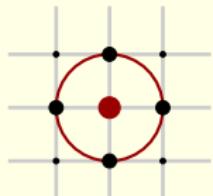
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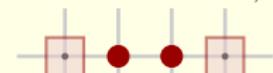
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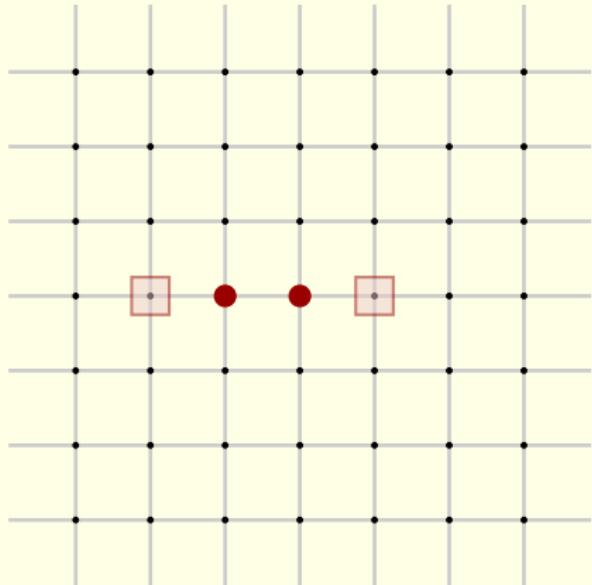
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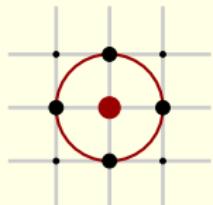
Horizontal pattern $\mathcal{S}_{r,\Delta}(\bullet \bullet)$:



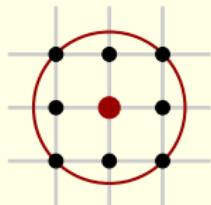
$$\Rightarrow D(r, \Delta) \geq 1/2$$



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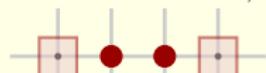


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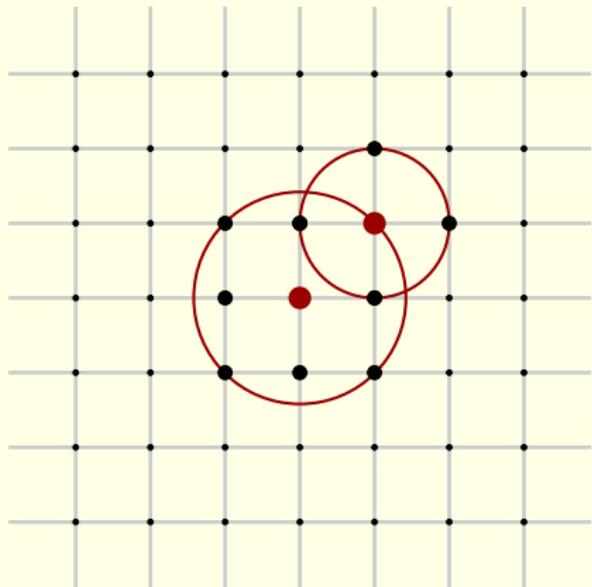


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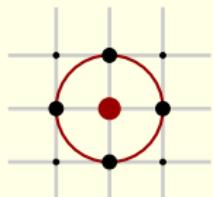
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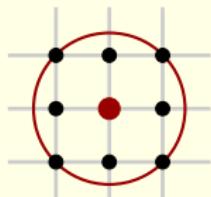
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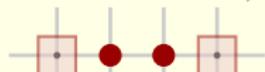


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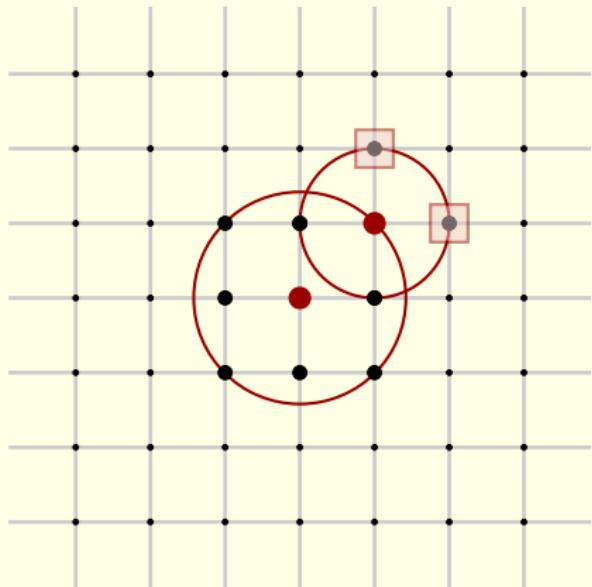


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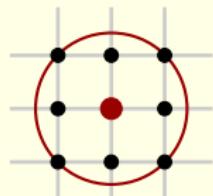
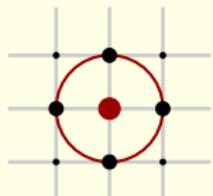
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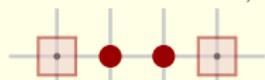
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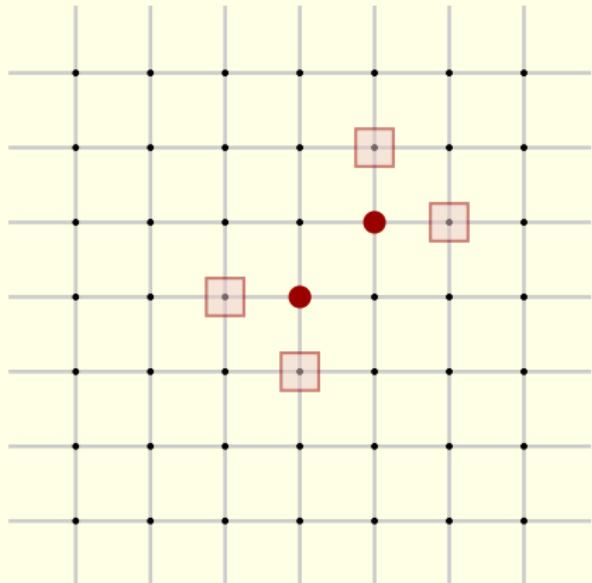
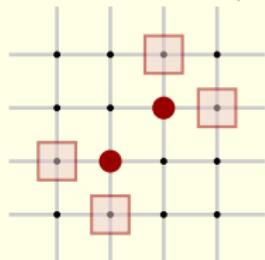


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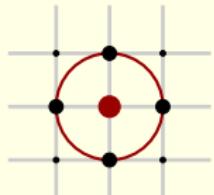


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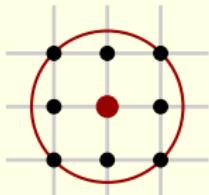
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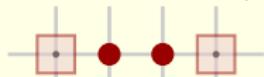


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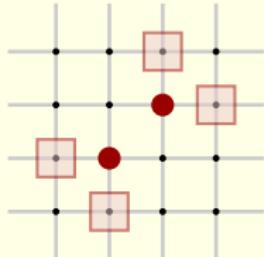
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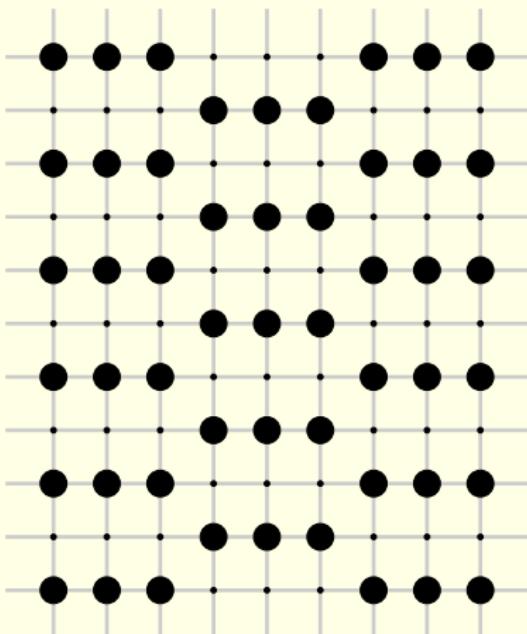


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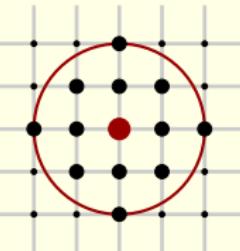
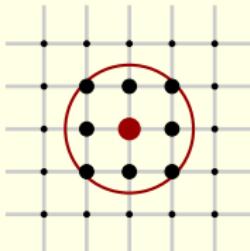
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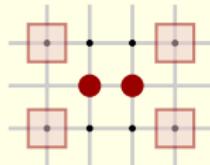
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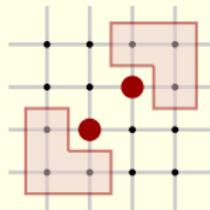
Next one: $(r, \Delta) = (\sqrt{2}, 2 - \sqrt{2})$



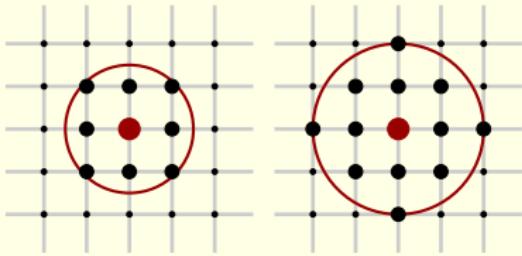
Horizontal pattern $\mathcal{S}_{r,\Delta}(\bullet \bullet)$:



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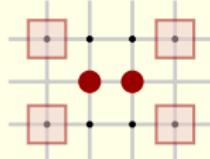


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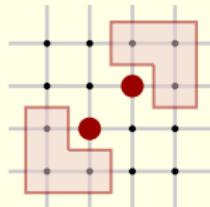


$$r = \sqrt{2} \quad r + \Delta = 2$$

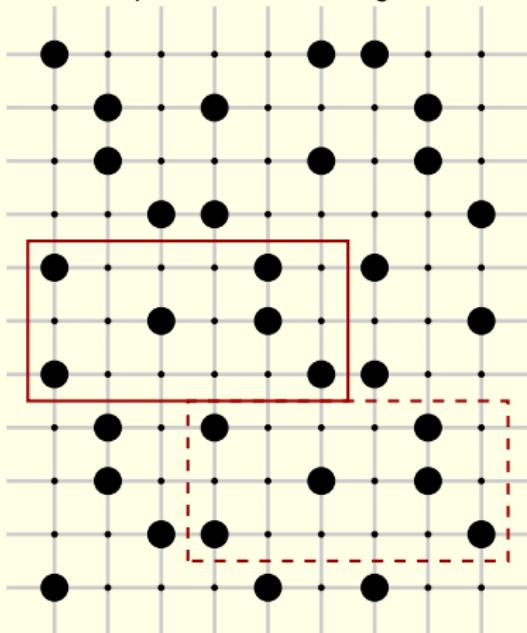
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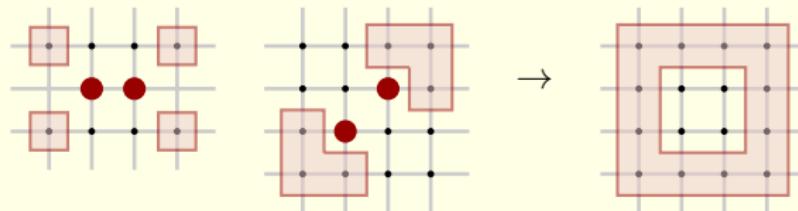
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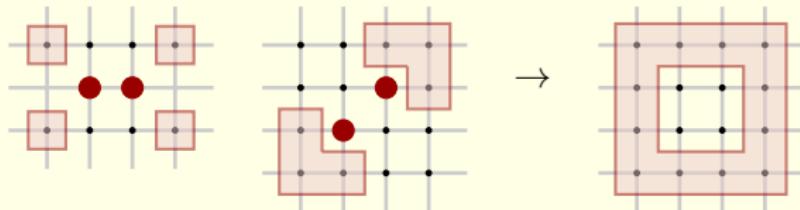
$$\frac{1}{4} \leq D(r, \Delta) \leq \frac{1}{3}:$$



Improvement of the lower bound

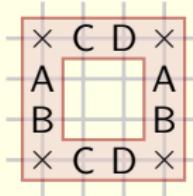


Improvement of the lower bound

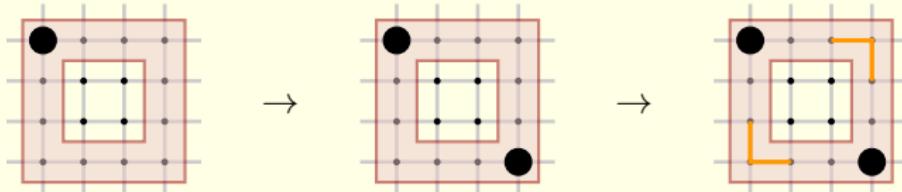


At least 3 code vertices in the frame:

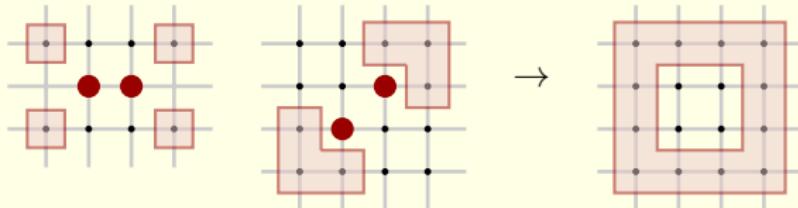
- If no code vertices in the corners, at least four codes in the frame:



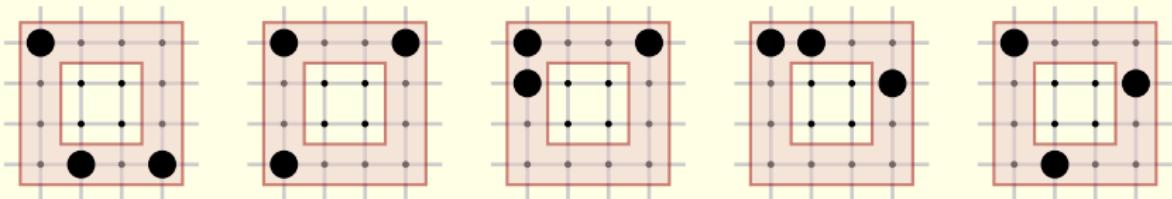
- if one code in a corner:



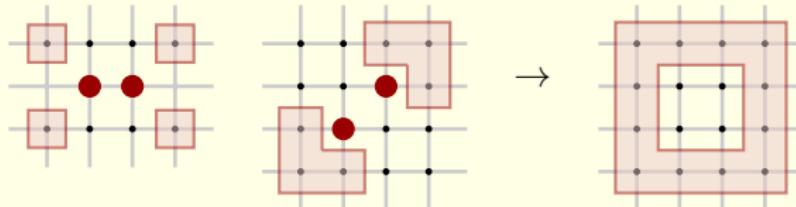
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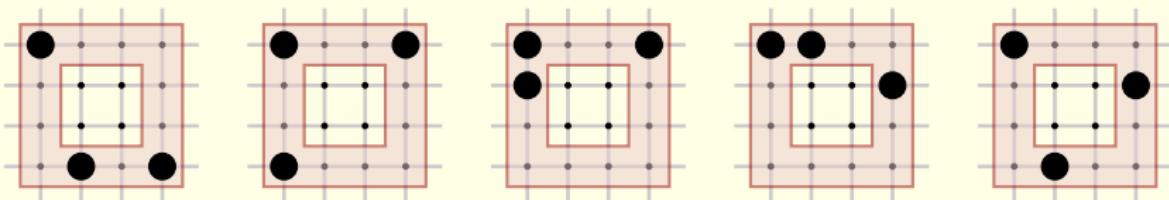
At least 3 codes vertices but only 5 possible configurations:



Improvement of the lower bound



At least 3 codes vertices but only 5 possible configurations:

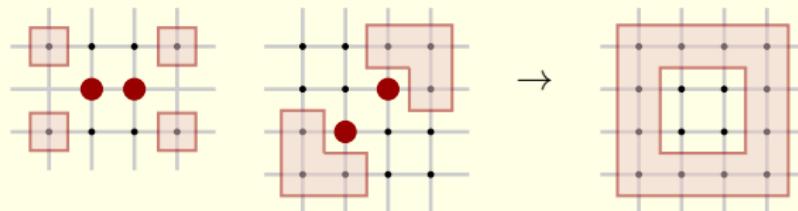


Each frame has a neighbor frame with at least 4 codes.

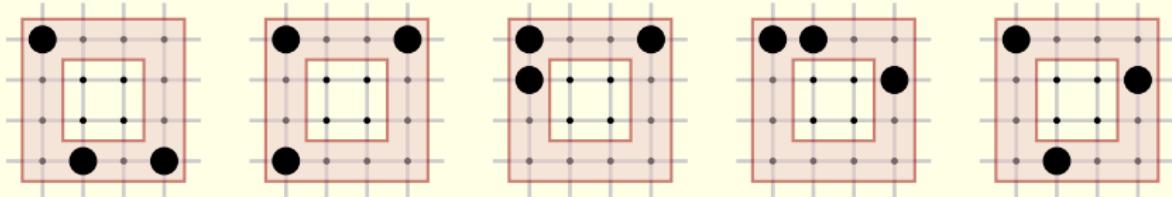
Each frame with at least 4 codes gives charge $1/5$ to neighbor frames with 3 codes.

$$D(\sqrt{2}, 2 - \sqrt{2}) \geq \frac{(3 + \frac{1}{5})}{12} = \frac{4}{15} \simeq 0.26$$

Improvement of the lower bound



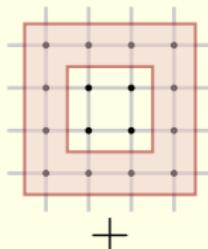
At least 3 codes vertices but only 5 possible configurations:



With other rules:

$$D(\sqrt{2}, 2 - \sqrt{2}) \geq \frac{16}{57} \simeq 0.28$$

A frame useful for other identification problems!



+

King Grid

+

(1, ≤ 2)-identifying code

=

Theorem (Laihonan, Foucaud & Parreau, 2011+)

Improvement of the lower bound from $\frac{5}{12}$ to $\frac{47}{111}$

(Upper Bound $\frac{3}{7}$)

Results for small values

| $r \setminus r + \Delta$ | 1 | $\sqrt{2}$ | 2 | $\sqrt{5}$ | $\sqrt{8}$ | 3 | $\sqrt{10}$ |
|--------------------------|---|-----------------|----------------------------------|-----------------|-----------------------------------|--------------------------------------|----------------------------------|
| 1 | 1 | $\frac{1}{2}^b$ | X | X | X | X | X |
| $\sqrt{2}$ | — | $\frac{2}{9}^a$ | $[\frac{16}{57}^c, \frac{1}{3}]$ | X | X | X | X |
| 2 | — | — | $[\frac{6}{35}, \frac{5}{29}]^a$ | $\frac{1}{2}^b$ | $\frac{1}{2}^b$ | X | X |
| $\sqrt{5}$ | — | — | — | $\frac{3}{8}^a$ | $[\frac{17}{100}^c, \frac{2}{9}]$ | $[\frac{1}{4}^b, \frac{1}{3}]$ | X |
| $\sqrt{8}$ | — | — | — | — | $\frac{1}{8}^a$ | $[\frac{1}{7}^c, \frac{4}{21}]$ | $[\frac{1}{4}^b, \frac{3}{8}]$ |
| 3 | — | — | — | — | — | $[\frac{1}{13.66}^a, \frac{1}{6}]^a$ | $[\frac{1}{6}^b, \frac{13}{48}]$ |

X: No (r, Δ) -identifying code

a: Previous results

b: Trivial lower bound with horizontal pattern

c: Lower bound with discharging rules

Remarks about Δ

To have existence: $\mathcal{S}_{r,\Delta}(\bullet \bullet) \neq \emptyset$

$$\Delta_m(r) = \sup\{\Delta | \mathcal{S}_{r,\Delta}(\bullet \bullet) \neq \emptyset\}$$

We have $\Delta_m(r) \leq 1$ et $\Delta_m(r) = 1$ iff $r \in \mathbf{N}$, and:

$$\Delta_m(r) \geq 1 - O\left(\frac{1}{\sqrt{r}}\right)$$

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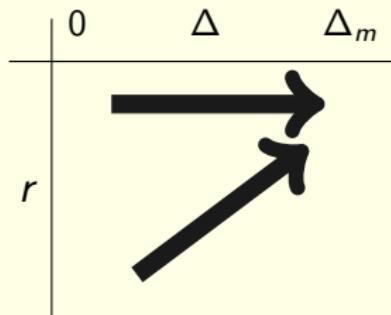
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Monotonicity of $D(r, \Delta)$:



General values of r

Lower bound: $D(r, \Delta) \geq \frac{1}{|\mathcal{S}_{r,\Delta}(\bullet \bullet)|} \rightarrow$ upper bound $|\mathcal{S}_{r,\Delta}(\bullet \bullet)|$

For big values of Δ : $D(r, \Delta) \geq \theta\left(\frac{1}{\sqrt{r}}\right)$

For other values: $D(r, \Delta) \geq D(r, 0) \geq \frac{1}{3.22r+4}$

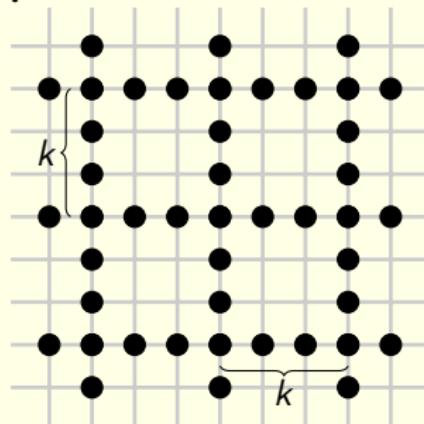
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Upper bound:



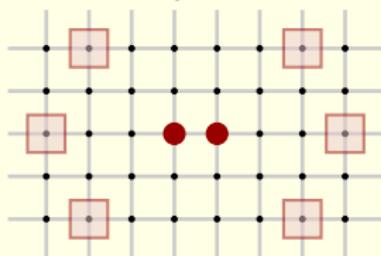
$$D(r, \Delta) \leq \frac{2}{k}$$

$$D(r, \Delta) \leq \frac{4}{r(2 - \Delta - \sqrt{2 - \Delta^2}) + C}$$

We can do better when we know the patterns

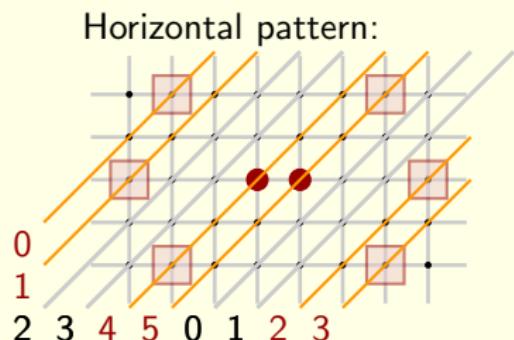
Example: $(r, \Delta) = (3, \sqrt{10} - 3)$:

Horizontal pattern:



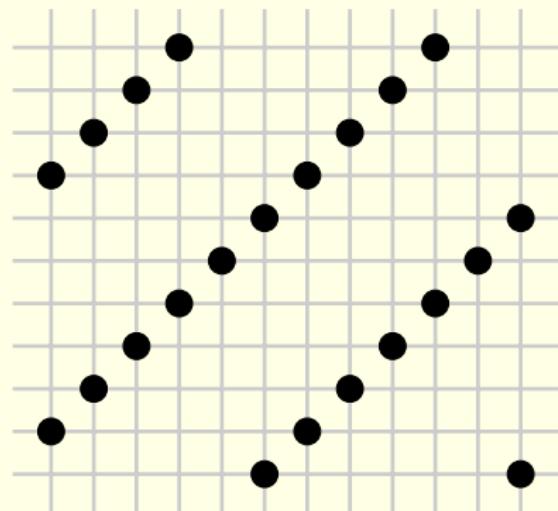
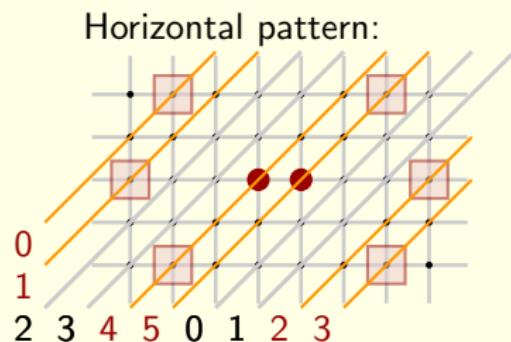
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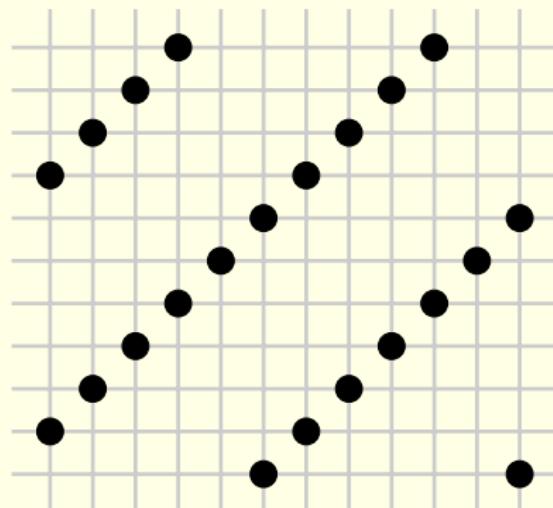
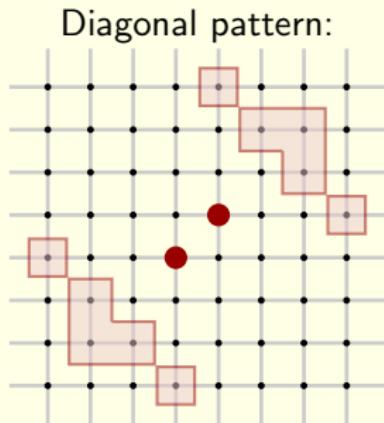
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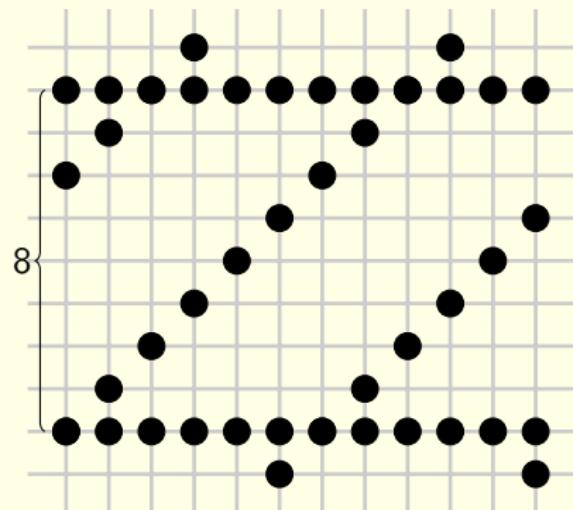
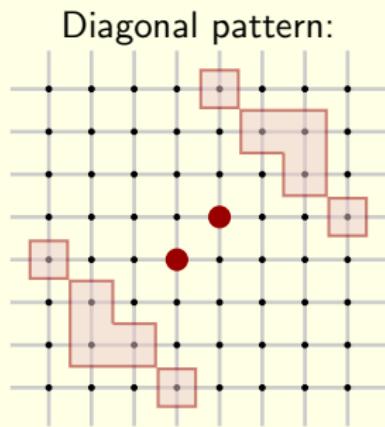
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Code of density $\frac{1}{6} + \frac{1}{8} - \frac{1}{48} = \frac{13}{48}$ (we can do $\frac{1}{4}$)

And even better

Sometimes, we can remove horizontal lines:

- $r = \sqrt{101}$, $r + \Delta = 11$, optimal density $\frac{1}{4}$
- $r = \sqrt{257}$, $r + \Delta = \sqrt{281}$, optimal density $\frac{1}{8}$

Why ?

- exactly d vertices in $\mathcal{S}_{r,\Delta}(\bullet \bullet)$ and all diagonals modulo d .
- d consecutive diagonals and anti-diagonals in $\mathcal{S}_{r,\Delta}(\bullet \bullet)$

⇒ Construction of infinite family with optimal codes

Example : $r = 2k^2 + 1$, $r + \Delta = \sqrt{r^2 + 2r - 3}$ with $k \not\equiv 0 [3]$, k odd: for r enough big, optimal code of density $\frac{1}{6}$

That's all

Thanks !